



Collecting telemetry data using P4 and RDMA

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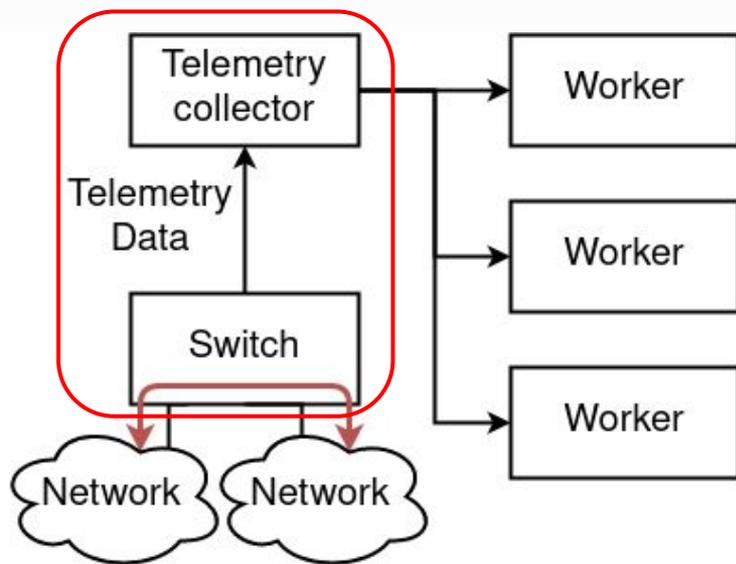
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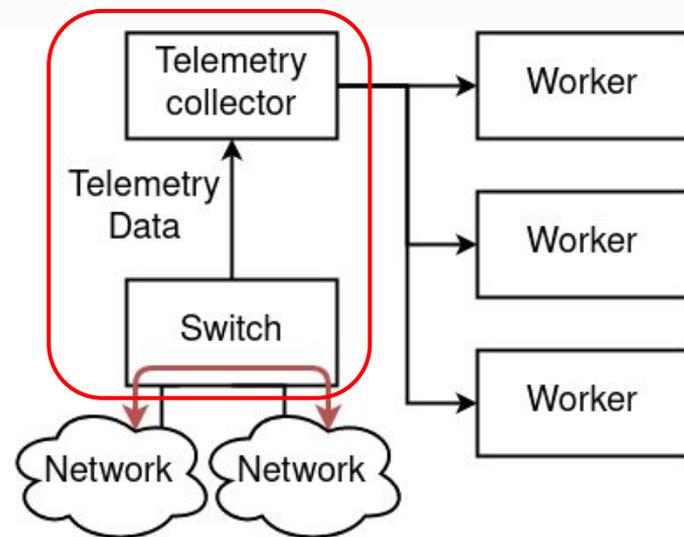
Introduction: Network Telemetry (I)

- ▶ Monitoring network health
- ▶ In-band network telemetry includes telemetry data in packets
- ▶ Delegate analyzation to multiple workers



Introduction: Network Telemetry (II)

- ▶ Requires an efficient means for collecting data
- ▶ Programming Protocol-independent Packet Processors (P4) for efficient telemetry data extraction
- ▶ Remote Direct Memory Access (RDMA) for efficient storage



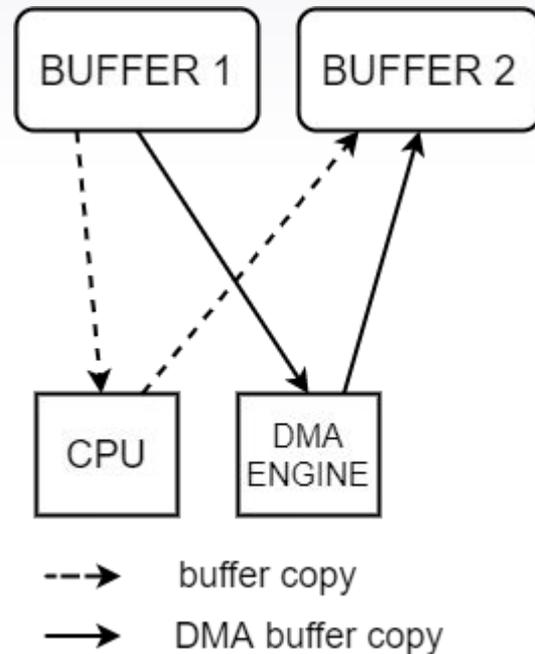
Research Questions

Can RDMA combined with P4 be used to efficiently collect telemetry data?

- ▶ How do we encapsulate telemetry data in an RDMA message?
- ▶ Can an RDMA session be maintained on a P4 switch?
- ▶ How can telemetry data be placed into persistent storage using RDMA?

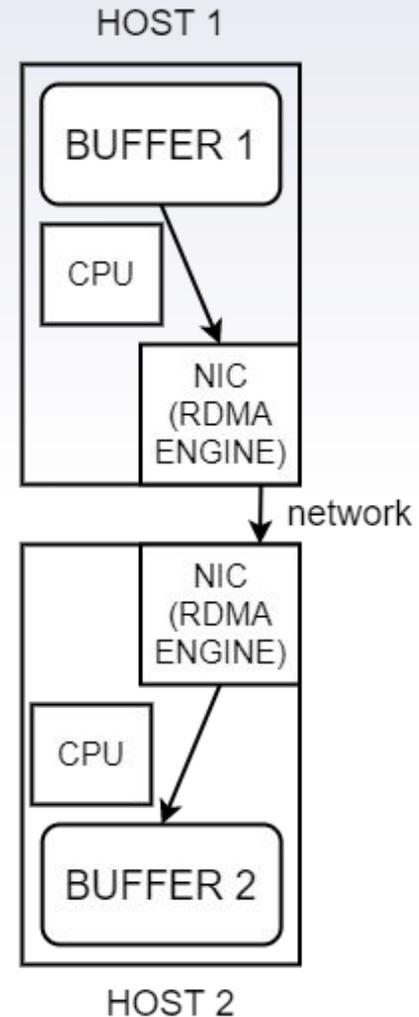
DMA

- ▶ Data is copied from buffer 1 to the buffer 2 via the CPU
- ▶ CPU spends a lot of cycles copying data
- ▶ Delegate high throughput transfers to DMA engine
- ▶ CPU can continue on other tasks while the DMA engine takes care of the transfer



RDMA

- ▶ Takes concept of DMA and puts it in the NIC
- ▶ Allows NIC to access data directly in memory
- ▶ CPU sets up a write operation
- ▶ The NIC on host 1 reads the buffer from memory and transfers it to the other NIC
- ▶ The NIC of host 2 writes the data to buffer 2
- ▶ The CPU is bypassed for the transfer of data



RoCEv1

- ▶ RDMA over Converged Ethernet version 1 (RoCEv1)
- ▶ RoCEv1 enables RDMA over layer 2 networks
- ▶ GRH has the same fields as IPv6
- ▶ BTH defines the RDMA operation for the NIC
- ▶ RETH includes memory address information for RDMA operations
- ▶ Invariant CRC is similar to Ethernet CRC, but slightly different



Related Work (I)

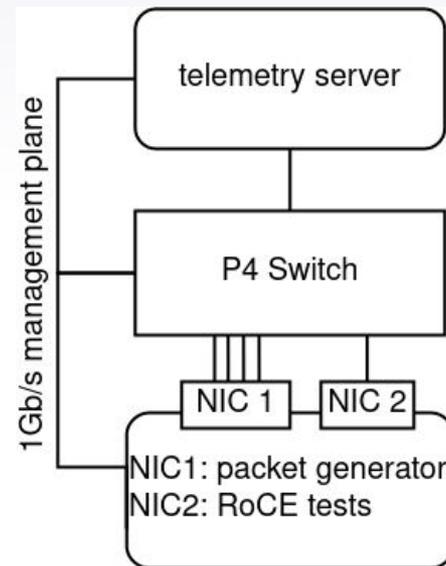
- ▶ Research by Tierney et al. (2012) compared the performance of TCP, UDP, UDT, and RoCE
 - ▶ CPU usage in RoCE is much less in comparison to the other protocols
 - ▶ RoCE showed consistently good performance
 - ▶ This research shows the potential of RoCE traffic in high-throughput networks

Related Work (II)

- ▶ Research by Kim et al. (2018) examined feasibility of implementing RoCE in P4 switch
 - ▷ Extending switch's buffer by storing burst data remotely
 - ▷ Extending forwarding tables by storing packet and action
 - ▷ Remotely increase counters for telemetry data
- ▶ “Borrowing” memory from remote server
- ▶ In our approach the server will eventually process this data further into the telemetry pipeline

Methodology & Setup

- ▶ Extract telemetry data with P4
- ▶ Implementing RoCE in P4 switch
- ▶ Send RoCE packet (RDMA write-only) with telemetry in payload
- ▶ Store payload on telemetry server

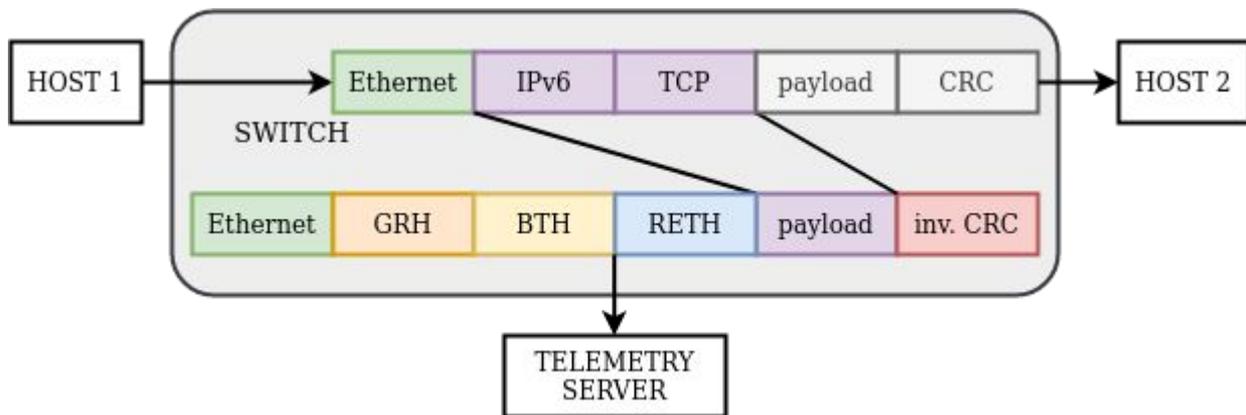


Server implementation

- ▶ Server uses *mmap* function to map virtual memory to a file on disk
- ▶ Set up the NIC to allow RDMA operations to the virtual memory address
- ▶ RDMA write-only can write directly to virtual memory, bypassing the CPU
- ▶ Open TCP socket to switch and share parameters required for RoCE packets

Switch implementation

- ▶ As there is no native support for RoCE on the switch, we create the RoCE headers from scratch in P4
- ▶ We learned the field values from the specification and experimentation



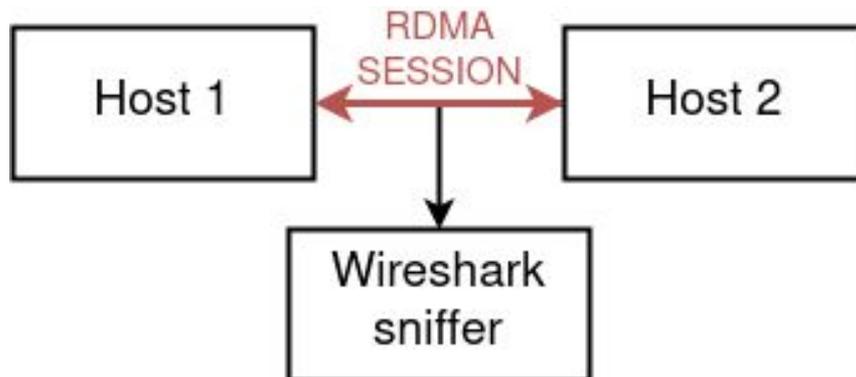
Switch: specific values

- ▶ Most of the header field values are static
- ▶ Others are dynamic or based on the server's RDMA parameters
 - ▷ Sequence number: counter increases with each packet
 - ▷ RDMA parameters from server are stored in a forwarding table
 - ▷ When the packet's egress port is to the telemetry server,
 - ▷ there is a match in the table
 - ▷ and the parameters are assigned to the packet
 - ▷ The virtual memory address is increased using an offset
 - ▷ CRC is calculated using an external function of the switch

Experiments (I)

Experiment 1: RoCEv1 experimentation to examine headers

- ▶ Establishing RDMA session between the two servers using RoCE libraries
- ▶ Analyze parameters that are used in the application and compare them to network traffic



Results experiment 1

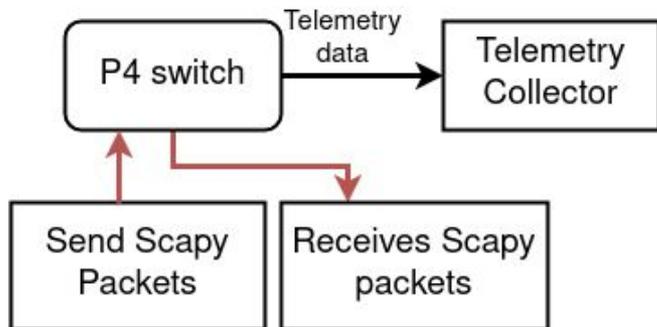
```
▼ InfiniBand
  ▼ Global Route Header
    0110 .... = IP Version: 6
    .... 0000 0000 .... = Traffic Class: 0
    .... .... 0000 0000 0000 0000 = Flow Label: 0
    Payload Length: 56
    Next Header: 27
    Hop Limit: 64
    Source GID: ::ffff:10.1.2.1
    Destination GID: ::ffff:10.1.2.2
  ▼ Base Transport Header
    Opcode: Reliable Connection (RC) - RDMA WRITE Only (10)
    0... .... = Solicited Event: False
    .1.. .... = MigReq: True
    ..11 .... = Pad Count: 3
    .... 0000 = Header Version: 0
    Partition Key: 65535
    Reserved: 00
    Destination Queue Pair: 0x000932
    1... .... = Acknowledge Request: True
    .000 0000 = Reserved (7 bits): 0
    Packet Sequence Number: 1
  ▼ RETH - RDMA Extended Transport Header
    Virtual Address: 26600384
    Remote Key: 178013
    DMA Length: 21
    Invariant CRC: 0xdd492f73
  ▼ Data (24 bytes)
    Data: 52444d41207772697465206f706572617469666e00000000
    [Length: 24]
```

```
[rutger@sne-dtn-04 rdma-writeonly]$ ./rdma-
tutorial -d mlx5_1 -i 1 -g 2
...
TCP connection was established
...
MR was registered with addr=0x195e3c0,
lkey=0x2b75d, rkey=0x2b75d, flags=0x7
QP was created, QP number=0x932
...
completion was found in CQ with status 0x0
Contents of server buffer: 'RDMA write
operation'
test result is 0
```

Experiments (II)

Experiment 2: RoCEv1 switch implementation testing

- ▶ Sending TCP packets crafted by Scapy from the Dell server
- ▶ Analyzed the file on the server to analyze correctness of the implementation



```
sendp(Ether()/IPv6(  
    dst="fc00::1111:2222:3333:4444",  
    src="fc00::5555:6666:7777:8888")/  
    TCP(dport=111, sport=222,  
        seq=0x1212, ack=0x3434),  
    iface="rename5")  
sendp(Ether()/IPv6(  
    dst="fc00::5555:6666:7777:8888",  
    src="fc00::1111:2222:3333:4444")/  
    TCP(dport=222, sport=111,  
        seq=0x3435, ack=0x1213),  
    iface="rename5")  
sendp(Ether()/IPv6(  
    dst="fc00::1111:2222:3333:4444",  
    src="fc00::5555:6666:7777:8888")/  
    TCP(dport=111, sport=222,  
        seq=0x1214, ack=0x3436),  
    iface="rename5")
```

Results experiment 2

```
Supermicro $ hexdump -C /mnt/nvme/output5
```

```
00 | fc00 0000 0000 0000 5555 6666 7777 8888  
10 | fc00 0000 0000 0000 1111 2222 3333 4444  
20 | 00de 006f 0000 1212 0000 3434 0000 0000  
30 | fc00 0000 0000 0000 1111 2222 3333 4444  
40 | fc00 0000 0000 0000 5555 6666 7777 8888  
50 | 006f 00de 0000 3435 0000 1213 0000 0001  
60 | fc00 0000 0000 0000 5555 6666 7777 8888  
70 | fc00 0000 0000 0000 1111 2222 3333 4444  
80 | 00de 006f 0000 1214 0000 3436 0000 0002
```

```
sendp(Ether()/IPv6(  
  dst="fc00::1111:2222:3333:4444",  
  src="fc00::5555:6666:7777:8888")/  
  TCP(dport=111, sport=222,  
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  TCP(dport=111, sport=222,  
      seq=0x1214, ack=0x3436),  
  iface="rename5")
```

Discussion

- ▶ No CPU involvement means CPU does not know anything about the data
- ▶ No signalling: signalling should provide method to let the CPU know when data can be read from memory
- ▶ P4 has no support for packet trailers, limiting the payload length

Conclusion

- ▶ RDMA is a feasible solution to communicate telemetry data to a collector
- ▶ P4 allows the original header to be encapsulated into a RoCE packet
- ▶ An RDMA session is maintained on the switch by keeping state of required parameters
- ▶ *mmap* provides the possibility of mapping a file to virtual memory, allowing RDMA access to this memory region

Future work

- ▶ Comparing the performance of this implementation with other techniques
 - ▷ Data Plane Development Kit (DPDK)
 - ▷ extended Berkeley Packet Filter (eBPF)
- ▶ Optimizing system performance (NVMe over Fabric instead of memory mapping)
- ▶ Investigate in an efficient method to signal the CPU that data can be processed further into the telemetry pipeline
 - ▷ RDMA write-only with immediate
- ▶ Completing the telemetry pipeline by adding workers

Security implications

- ▶ Remote key is equivalent to a plain text password
- ▶ According to RFC 5040 manufacturers MUST ensure that only memory in a specific Protection Domain can be accessed.
- ▶ Full security considerations in RFC 5040 and RFC 5042
- ▶ Throwhammer is an RDMA variant on the Rowhammer attack
- ▶ If properly set up, security implications similar to UDP/TCP streams (traffic injection/sniffing).

CRC calculation

```
CRCPolynomial<bit<32>>(
    coeff = 0x04C11DB7,
    reversed = true,
    msb = false,
    extended = false,
    init = 0xFFFFFFFF,
    xor = 0xFFFFFFFF) poly;
```

Local Route Header	0xFFFFFFFF
Global Route Header	...
	Traffic Class: 0xFF
	Flow Label: 0FFFFFF
	...
	Hop Limit: 0xFF
Base Transport Header	...
	Reserved: 0xFF
	...
RDMA Extended Transport Header	...
Payload	...

References

- ▶ (R)DMA figures inspired on:
http://www.rdmaconsortium.org/home/The_Case_for_RDMA020531.pdf