

DEVELOPMENT OF A NEW POLICY EVALUATION PROCEDURE FOR XACML

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WHY

- Customer data more and more valuable
- Data stored in cloud
- Access control becomes critical

XACML

- eXtensible Access Control Markup Language
- XML-based language
- Also an architecture
- OASIS standard for the expression of security policies

XACML ELEMENTS

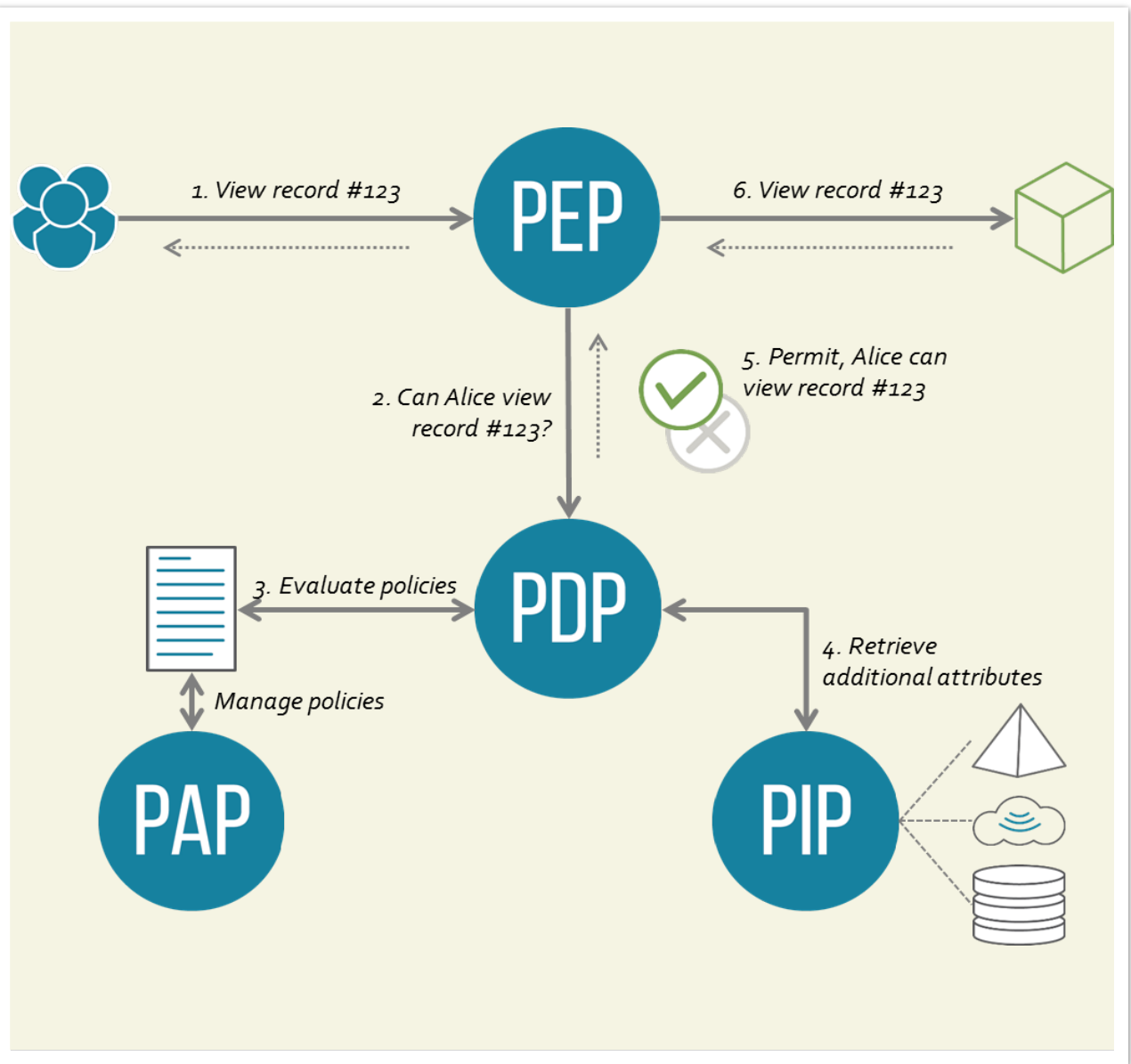
```
<PolicySet>
  <Policy RuleCombiningAlg="..." >
    <Target/>
    <Rule RuleId="..." Effect="Permit">
      <Target/>
      <Condition/>
    </Rule>
    <Rule RuleId="..." Effect="Deny">
      <Target/>
      <Condition/>
    </Rule>
  </Policy>
  <Policy RuleCombiningAlgId="...">
    ...
  </Policy>
</PolicySet>
```

EXAMPLE XACML POLICY

```
<Policy
  RuleCombiningAlgId="identifier:rule-combining-algorithm:permit-overrides">
  <Target/>
  <Rule RuleId="urn:oasis:names:tc:xacml:3.0:example:SimpleRule1" Effect="Permit">
    <Target>
      <AnyOf>
        <AllOf>
          <Match MatchId="string-equal">
            <AttributeValue DataType="string">admin</AttributeValue>
            <AttributeDesignator AttributeId="role" DataType="string"/>
          </Match>
        </AllOf>
      </AnyOf>
    </Target>
    <Condition>
      ...
    </Condition>
  </Rule>
  <Rule RuleId="urn:oasis:names:tc:xacml:3.0:example:SimpleRule2" Effect="Deny">
    ...
  </Rule>
</Policy>
```

XACML IN ACTION

1. Request intercepted by PEP
2. Request converted to XACML
3. PDP evaluates policy
4. If needed retrieve additional attributes
5. PDP reaches decision and forwards this to PEP
6. Request arrives at resource



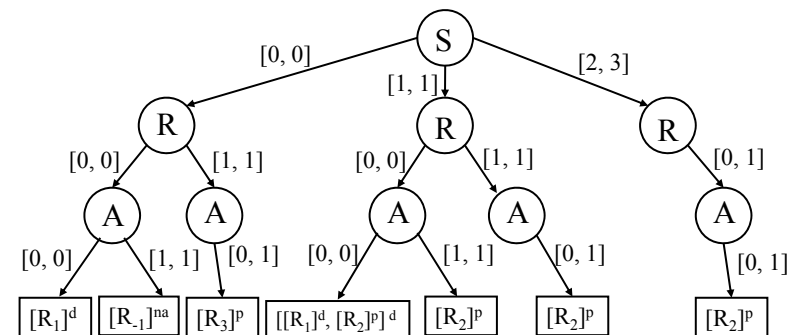
RELATED RESEARCH

(Adaptive) reordering

- Based on statistics and categorization

Decision Diagrams

- XEngine
- Matching Tree (MT) and Combining Tree (CT)
- SNE-XACML with MIDD



Source: XEngine: A Fast and Scalable XACML Policy Evaluation Engine

RESEARCH QUESTION

- Propositional encoding
- PDP

SAT & CNF

- Boolean function: $f(x_1, x_2, \dots, x_n)$
- Variables, operators and parentheses: $x_1, \wedge, \vee, \neg, ()$
- SAT solvers
- CNF: $(p_1 \vee p_2) \wedge (p_3 \vee p_4) \wedge (p_5 \vee p_6)$
 - Conjunction of clauses
 - Disjunction of literals

ALGORITHM

- Constructing attribute domains
- Policy flattening
- SAT encoding

CONSTRUCTING ATTRIBUTE DOMAINS (1)

- Attributes
 - *AttributeValue*
 - *AttributeDesignator*
 - *AttributeSelector*

```
<rule Effect="Permit">
```

```
...
```

```
<AttributeValue DataType="String">admin</AttributeValue>
```

```
<AttributeDesignator AttributeId="role" DataType="String"/>
```

```
...
```

```
</rule>
```

CONSTRUCTING ATTRIBUTE DOMAINS (2)

$$D_{role} \in \{\text{admin, manager, hr, user}\}$$

$$\text{admin} \in \{\text{admin, manager, hr, user}\}$$

CONSTRUCTING ATTRIBUTE DOMAINS (3)

Algorithm 1 EnumerateVariables

Input: A map m containing the DataTypes as keys and (empty) arrays as values and a policy p

```
1: procedure ENUMERATEVARS( $p, m$ )
2:   for all target elements do
3:     update  $m$  with values found in the policy
      target
4:   end for
5:   for all variable definitions do
6:     update  $m$  with values found in the variable
      definitions
7:   end for
8:   for all policy elements do
9:     if element is a policy then
10:      enumerateVars(element,  $m$ )
11:    else if element is a rule then
12:      update  $m$  with values found in the rule
        targets
13:      update  $m$  with values found in the rule
        condition
14:    end if
15:  end for
16: end procedure
```

ALGORITHM

- Constructing attribute domains
- Policy flattening
- SAT encoding

Policy Flattening (1)

- Applicability space $\langle AS_A, AS_{IN}, AS_{NA} \rangle$
- Decision space $\langle DS_P, DS_D, DS_{IN}, DS_{NA} \rangle$

Algorithm 2 FlattenPolicy

Input: A policy p **Output:** Decision space

$< DS_P, DS_D, DS_{IN(P)}, DS_{IN(D)}, DS_{IN(NA)}, DS_{NA} >$

```
1: procedure FLATTENPOLICY( $p$ )
2:   if  $p$  is a rule then
3:      $AS_A^P = AS_A^T \cap AS_A^C$ 
4:      $AS_{IN}^P = AS_{IN}^C \cup AS_{IN}^T$ 
5:     if effect of  $p$  is Permit then
6:        $DS_P = AS_A^P$ 
7:        $DS_D = \emptyset$ 
8:        $DS_{IN(P)} = AS_{IN}^P$ 
9:        $DS_{IN(D)} = \emptyset$ 
10:    else if effect of  $p$  is Deny then
11:       $DS_P = \emptyset$ 
12:       $DS_D = AS_A^P$ 
13:       $DS_{IN(P)} = \emptyset$ 
14:       $DS_{IN(D)} = AS_{IN}^P$ 
15:    end if
16:     $DS_{IN(PD)} = \emptyset$ 
17:     $DS_{IN(NA)} =$ 


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     $(DS_P \cup DS_D \cup DS_{IN(P)} \cup DS_{IN(D)} \cup DS_{IN(PD)})$ 
18:    return
     $(DS_P, DS_D, DS_{IN(P)}, DS_{IN(D)}, DS_{IN(PD)}, DS_{IN(NA)})$ 
19:  else if  $p$  is a policy (set) then
20:    policies =  $\emptyset$ 
21:    for all elements  $e$  of  $p$  do
22:      result = flattenPolicy( $e$ )
23:      add result to policies
24:    end for
25:    combiningAlg = combining algorithm of  $p$ 
26:    return applyCA(policies, combiningAlg)
27:  end if
28: end procedure
```

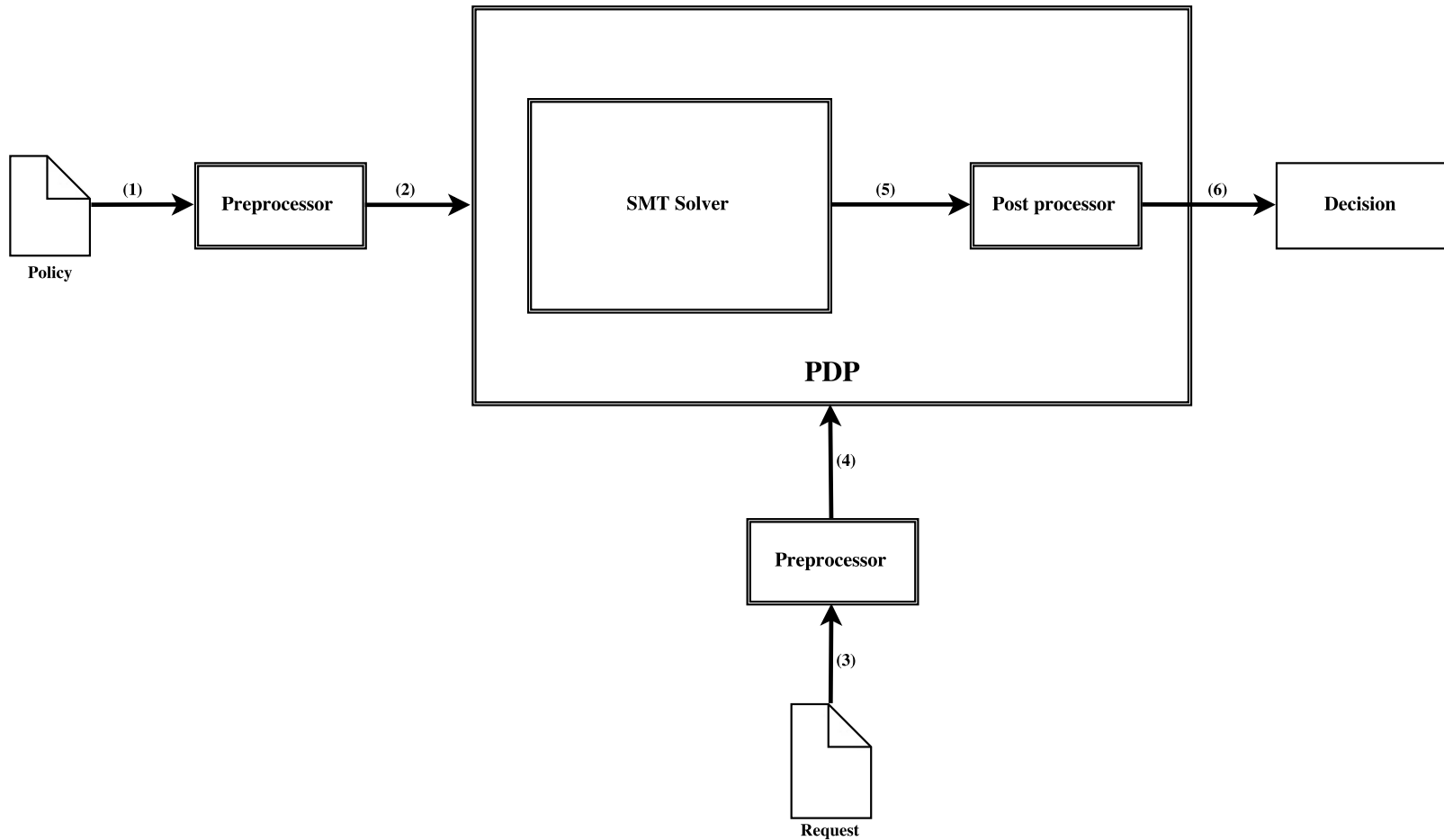
ALGORITHM

- Constructing attribute domains
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SAT ENCODING

$$DS_P \cup DS_D \cup DS_{IN(P)} \cup DS_{IN(D)} \cup DS_{IN(PD)} \cup DS_{NA}$$

FRAMEWORK



Conclusion

- Creating SAT formula
- SAT solvers
- No trees
- Experimental validation